

Mid Term

- This mid-term is due on **April 5, 2021**.
- Each problem carries 10 points.
- No collaboration is allowed for mid-term!

1. **(Infinite Integrals)** In the following, let $f : [0, 1] \rightarrow \mathbb{R}$ be continuous. **Rigorously** evaluate the following limits.

$$(1) \quad \lim_{n \rightarrow \infty} \int_0^1 \int_0^1 \dots \int_0^1 f\left(\frac{x_1 + x_2 + \dots + x_n}{n}\right) dx_1 dx_2 \dots dx_n$$

$$(2) \quad \lim_{n \rightarrow \infty} \int_0^1 \int_0^1 \dots \int_0^1 f\left((x_1 x_2 \dots x_n)^{1/n}\right) dx_1 dx_2 \dots dx_n.$$

$$(3) \quad \lim_{n \rightarrow \infty} \int_0^1 \int_0^1 \dots \int_0^1 \frac{x_1^2 + x_2^2 + \dots + x_n^2}{x_1 + x_2 + \dots + x_n} dx_1 dx_2 \dots dx_n$$

2. **(Superharmonic functions)** A function $\phi : \mathbb{Z}^k \rightarrow \mathbb{R}$ on the integer lattice \mathbb{Z}^k is called *superharmonic* if for each lattice point \mathbf{x} , we have

$$\phi(\mathbf{x}) \geq \frac{1}{2k} \sum_{\mathbf{y}: \|\mathbf{y}-\mathbf{x}\|_1=1} \phi(\mathbf{y}),$$

where the sum extends over the $2k$ nearest neighbors. Using Martingale methods, show for $k = 1$ and $k = 2$ that all bounded superharmonic functions are constant.

HINT: You may use the fact that the symmetric random walk on \mathbb{Z}^k is recurrent for $k = 1, 2$.

3. **(Find the odd one out)** Assume that you have m coins. All the coins are unbiased (that is, they have an equal probability of heads or tails), EXCEPT the *special* coin which comes up heads with probability $\frac{1}{2} + \rho$ for some unknown $0 < \rho < \frac{1}{2}$. Your goal is to identify the special coin. For this, you toss each of the coins n times and record the results of each toss. Come up with a decision rule to infer the special coin from the experimental results. Find the sample complexity (*i.e.*, a lower bound to n) of your decision rule to ensure that the probability of error is at most δ .
4. **(Gaussian Complexity of ℓ_0 -“balls”)** Sparsity plays an important role in many high-dimensional statistical models. The Gaussian complexity of an Euclidean set $S \subseteq \mathbb{R}^d$ is defined as:

$$\mathcal{G}(S) = \mathbb{E}\left(\max_{\mathbf{x} \in S} \sum_{i=1}^d x_i Z_i\right),$$

where $\{Z_i, 1 \leq i \leq d\}$ are i.i.d. standard Gaussian random variables. In this problem, we will compute the Gaussian complexity of an s -sparse ℓ_0 -ball intersected with a unit ℓ_2 -ball. Consider the set

$$T^d(s) = \{\theta \in \mathbb{R}^d \mid \|\theta\|_0 \leq s, \|\theta\|_2 \leq 1.\}$$

corresponding to all s -sparse vectors contained within the Euclidean unit ball. In this problem, we prove that its Gaussian complexity is upper bounded as:

$$\mathcal{G}(T^d(s)) \lesssim \sqrt{s \log \left(\frac{ed}{s} \right)}. \quad (1)$$

- (a) First show that $\mathcal{G}(T^d(s)) = \mathbb{E}[\max_{|S|=s} \|w_S\|_2]$, where $w_S \in \mathbb{R}^{|S|}$ denotes the sub-vector of (w_1, w_2, \dots, w_d) indexed by the subset $S \subseteq \{1, 2, \dots, d\}$.
- (b) Next show that for any fixed subset S of cardinality s :

$$\mathbb{P}[\|w_S\|_2 \geq \sqrt{s} + \delta] \leq e^{-\delta^2/2}.$$

- (c) Use the preceding parts to establish the bound (1).